Probabilistic Graphical Models

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1 The treewidth of a graph

A crucial property to identify if probabilitic inference in a graphical model can be done efficiently is the *treewidth* of the underlying graph.

Definition 1. (see Wikipedia: https://en.wikipedia.org/wiki/Chordal_graph) A graph is called *chordal*, if any cycle in it that consists of four or more vertices has a *chord*, i.e. there exists an edge that is not part of the cycle but connects two vertices of the cycle.

Definition 2. (see Wikipedia: https://en.wikipedia.org/wiki/Chordal_completion) A chordal completion of a graph is a chordal graph that has the same vertex set and contains at least all edges of the original graph. Note: in general, graphs can have many different chordal completions.

Definition 3. (see Wikipedia: https://en.wikipedia.org/wiki/Treewidth) The *treewidth* of a chordal graph is the size of its largest clique minus 1. The *treewidth* of a (potentially non-chordal) graph is the smallest treewidth of any of its chordal completions.

For each the following graphs 1)-7),

- a) determine if it is chordal,
- b) if not, construct a chordal completion,
- c) determine its treewidth.





a) graph that is not chordal

b) not a chordal completion of a)





c) a chordal completion of a)

d) a chordal completion of a)



e) a chordal completion of a)

f) chordal, but not a completion of a)





 f_2



Hint for 1h): ... chord on and chord of a chord.

2 Factor Graphs

Assume you are given eight binary-valued random variables, X_1, \ldots, X_8 . Construct factor graphs for the following probability distributions (with $x = (x_1, \ldots, x_8)$), such that their underlying graphs have minimal treewidth.

- a) $p(x) \propto e^{\text{number of 1s in } x}$
- b) $p(x) \propto e^{\text{number of } (0-1) \text{ transitions in } x_1, \dots, x_8}$
- c) $p(x) \propto e^{\text{number of } (0-1-0) \text{ transitions in } x_1, \dots, x_8}$
- d) $p(x) \propto e^{\text{number of (1-0-1) combinations between any three distinct entries in } x}$

e)
$$p(x) = \begin{cases} 1 & \text{if } x = (0, 0, \dots, 0) \\ 0 & \text{otherwise} \end{cases}$$

f)
$$p(x) = \begin{cases} \frac{1}{2} & \text{if } x = (1, 1, \dots, 1) \\ \frac{1}{510} & \text{otherwise} \end{cases}$$

g)
$$p(x) \propto e^{\text{parity of } x}$$

Could you do better, if you introduced additional (latent) random variables?

3 Marginal Inference

Assume you are given four binary-valued random variables, X_1, \ldots, X_4 , and a distribution $p(x_1, x_2, x_3, x_4) \propto \phi_1(x_1, x_2)\phi_2(x_2, x_3)\phi_3(x_3, x_4)$ with factors $\phi_i(x_i, x_{i+1}) = \begin{cases} 3 & \text{if } x_i = 0 \text{ and } x_{i+1} = 1 \\ 1 & \text{otherwise} \end{cases}$ for $i = 1, \ldots, 3$.

Compute (on paper!):

- a) the normalizing constant
- b) the probability $p(x_1 = 0, x_2 = 0, x_3 = 0, x_4 = 0)$
- c) the marginal probability $p(x_1 = 0)$
- d) $\operatorname{corr}(X_1, X_2)$
- e) the marginal probability $p(x_1 = 0, x_4 = 0)$

In each case, perform the computation in two ways: once naively, and once using belief propagation where possible (note: e) might require some thought for this). What is more efficient? How would this change f) for a larger number of variables, g) for variables with more states?

4 Maximum Entropy Distribution

Complete the proof we skipped in the lecture:

For samples z^1, \ldots, z^N and feature functions $\phi_i : \mathbb{Z} \to \mathbb{R}$ for $i = 1, \ldots, d$, define $\mu_i := \sum_{n=1}^N \phi_i(z_i)$. Show for finite \mathbb{Z} : out of all probability distribution, p(z), that fulfill $\mathbb{E}_{z \sim p(z)}[\phi(z)] = \mu_i$ for $i = 1, \ldots, D$, the one with *highest entropy* has the form

$$p(z) \propto \exp\left(\sum_{i} \theta_{i} \phi_{i}(z)\right)$$
 for some values $\theta_{1}, \dots, \theta_{D} \in \mathbb{R}$.